

# OVERVIEW OF "ALGORITHMIC LEARNING THEORY AND CRYPTOGRAPHY"

Tobias Eibach

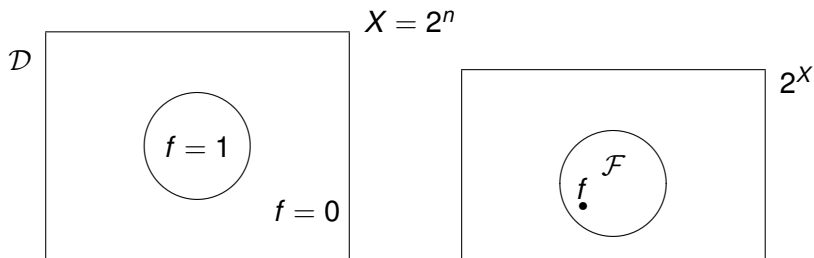
Institute of Theoretical Computer Science  
University of Ulm

25.4.2007

# OUTLINE

- 1** INTRODUCTION TO ALGORITHMIC LEARNING THEORY
- 2** IMPACTS ON CRYPTOGRAPHY
- 3** LEARNING VS CRYPTOGRAPHY
- 4** REFERENCES / RESEARCH TOPICS

# PAC LEARNING SCHEMA



Learner  
knows:  $\mathcal{D}, \mathcal{F}$

Teacher  
knows:  $\mathcal{D}, f \in \mathcal{F}$

request  
example  $\langle x_i, f(x_i) \rangle$

# INTRODUCTION TO ALGORITHMIC LEARNING THEORY

## DEFINITION

An algorithm  $A$  learns a class of functions  $\mathcal{F}$  if  $\forall f \in \mathcal{F}$  and  $\epsilon, \delta > 0$ , the algorithm  $A$  outputs an hypothesis  $h$  with probability  $1 - \delta$  such that

$$\text{error}(f, h) \leq \epsilon$$

$$\text{error}(f, h) := \Pr_{x \in \mathcal{D}}[f(x) \neq h(x)]$$

The running time is polynomial if it's polynomial in  $n$ ,  $1/\epsilon$  and  $\log(1/\delta)$ .

# MAIN LEMMA

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Let  $f$  be a Boolean function on  $n$  variables computable by a Boolean circuit of depth  $d$  and size  $M$ , and let  $t$  be any integer, then

$$\sum_{S \subseteq \{1..n\}, |S| \geq t} \hat{f}(S)^2 \leq M 2^{-t^{1/d}/20}$$

By N. Lineal, Y. Mansour and N. Nisan.

# MAIN LEMMA

The Main Lemma follows from the following 2 results:

## LEMMA (HASTAD)

$$\Pr_{\rho}[\text{DT-depth}(f_{\rho}) \geq s] \leq M2^{-s}$$

Where  $\rho$  is a random restriction with parameter  $p \leq \frac{1}{10^d s^{d-1}}$ .

## LEMMA

$$\sum_{|S|>t} \hat{f}^2(S) \leq 2\Pr_{\rho}[\text{DT-depth}(f_{\rho}) \geq tp/2]$$

## COROLLARY

Functions in  $AC^0$  can be learned efficiently.

# NO PSEUDORANDOM FUNCTION GENERATORS IN $AC^0$

$f : \{0, 1\}^m \times \{0, 1\}^n \rightarrow \{0, 1\}$  is called PRFG if no oracle TM  $M$  running in polynomial time can distinguish between a truly random oracle and the oracle  $f(s, *)$ ,  $s$  chosen at random.

## COROLLARY

There exists no PRFG in  $AC^0$ .



# A PRIVATE KEY CRYPTOSYSTEM

A private cryptosystem based on a non-learnable function class (on the average)

- natural mapping
- private key CS  $(G, E, D)$
- first  $G$  generates a function  $f$  represented by  $\sigma$
- $D(E(m, \sigma), \sigma) = m$
- encrypt 0 by a neg. example and 1 by a pos. example

# REFERENCES



N. Lineal, Y. Mansour, N. Nisan.

*Constant Depth Circuits, Fourier Transform, and Learnability.*

Journal of the ACM, Vol. 40, No. 3, 1993, pp. 607-620.



A. Blum, M. Furst, M. Kearns, R. Lipton.

*Cryptographic Primitives Based on Hard Learning Problems.*

Lecture Notes in Computer Science, Vol. 773, 1994, pp. 278-291.

# RESEARCH TOPICS

- Lower Bounds on Cryptographic Primitives
- Relationship between A.L. and Zero Knowledge
- Applications of Game Theory in Cryptography and Algorithmic Learning Theory
- Latest developements in A.L.T. (learning of juntas, sensitivity of monotone decision trees, noisetolerance learning, agnostic learning) and their influence in Cryptography.
- Implementation of Steganographic Tools

# THE END

Thank you!

(full talk: <http://theorie.informatik.uni-ulm.de/Personen/eibach/> )